



# **Motion Planning**

METR 4202: Robotics & Automation

Dr Surya Singh -- Lecture # 10

October 5, 2016

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http://robotics.itee.uq.edu.au/~metr4202/

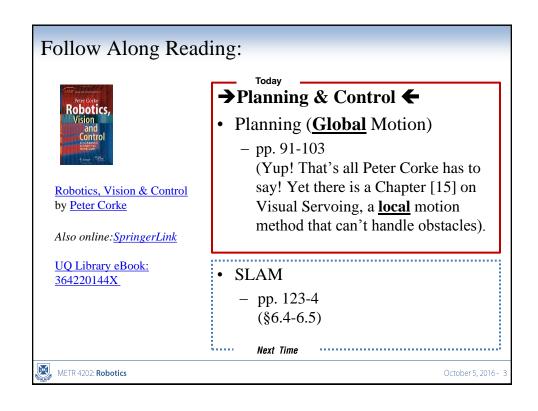
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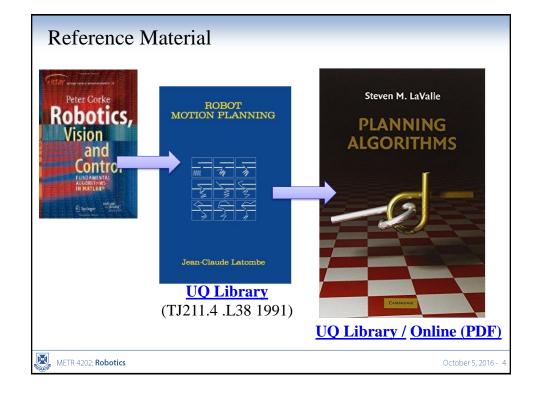
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### Schedule of Events

Week	Date	Lecture (W: 12:05-1:50, 50-N202)
1	27-Jul	Introduction
2	3-Aug	Representing Position & Orientation & State
		(Frames, Transformation Matrices & Affine Transformations)
3	10-Aug	Robot Kinematics Review (& Ekka Day)
4	17-Aug	Robot Inverse Kinematics & Kinetics
5	24-Aug	Robot Dynamics (Jacobeans)
6	31-Aug	Robot Sensing: Perception & Linear Observers
7	7-Sep	Robot Sensing: Single View Geometry & Lines
8	14-Sep	Robot Sensing: Feature Detection
9	21-Sep	Robot Sensing: Multiple View Geometry
	28-Sep	Study break
10	5-Oct	Motion Planning
11	12-Oct	Probabilistic Robotics: Localization & SLAM
12	19-Oct	Probabilistic Robotics: Planning & Control
13	26-Oct	State-Space Automation (Shaping the Dynamic Response/LQR) + Course Review

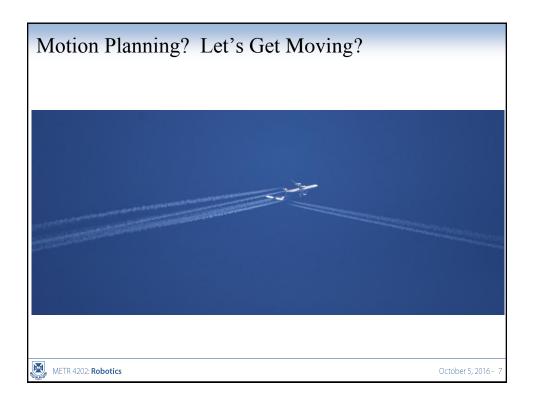
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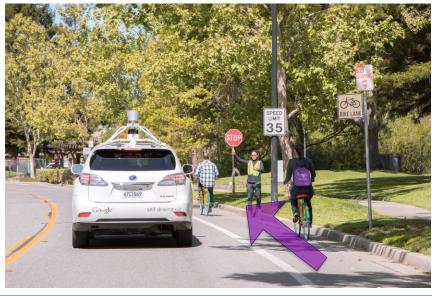
# (Kinematic) Motion Planning







### Motion Planning: Processing the Limits



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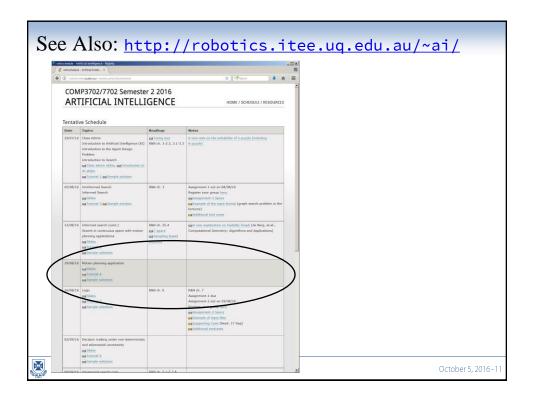
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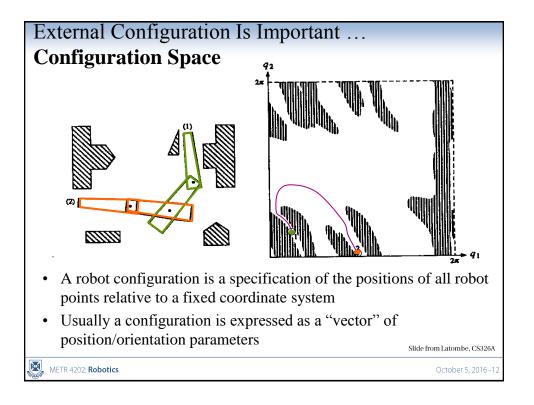
# Path-Planning Approaches

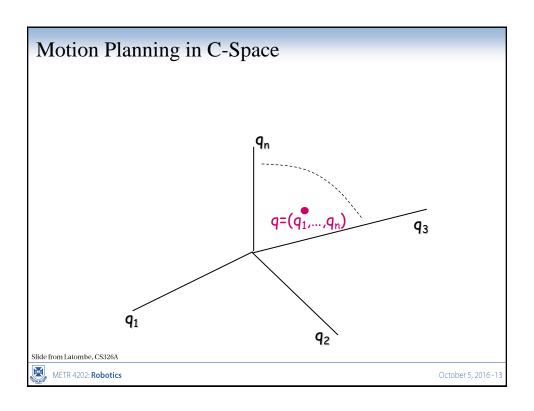
- Roadmap
   Represent the connectivity of the free space by a network of 1-D curves
- Cell decomposition
   Decompose the free space into simple cells and represent the connectivity of the free space by the adjacency graph of these cells
- Potential field
   Define a function over the free space that has a global minimum at the goal configuration and follow its steepest descent

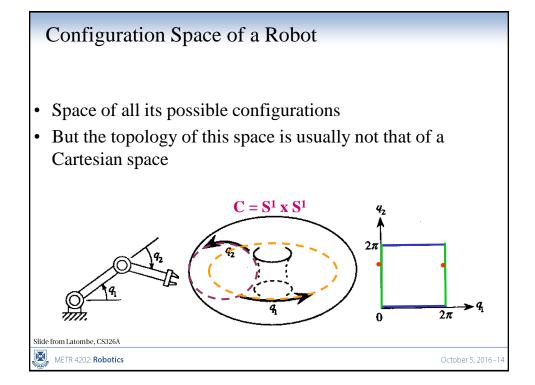
Slide from Latombe, CS326A

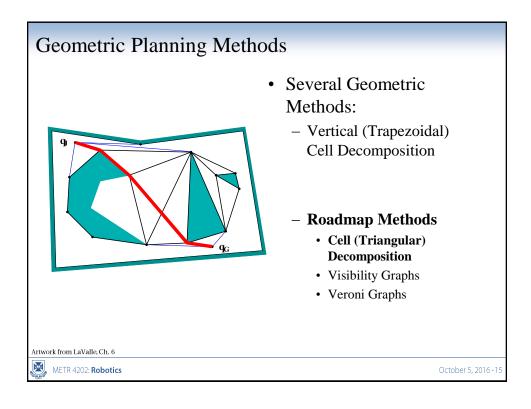


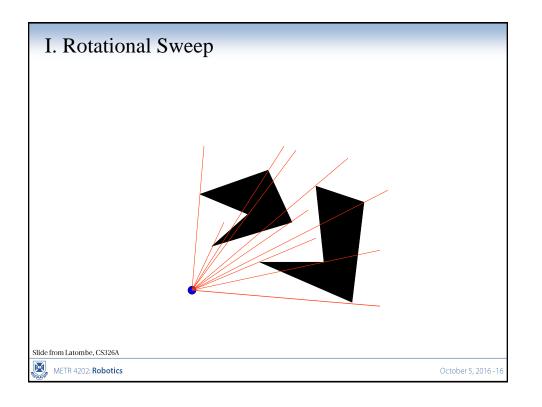


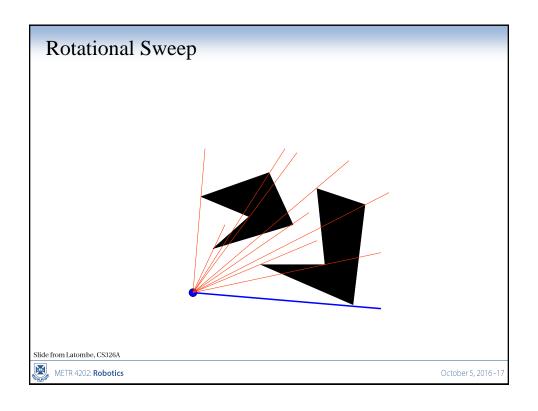


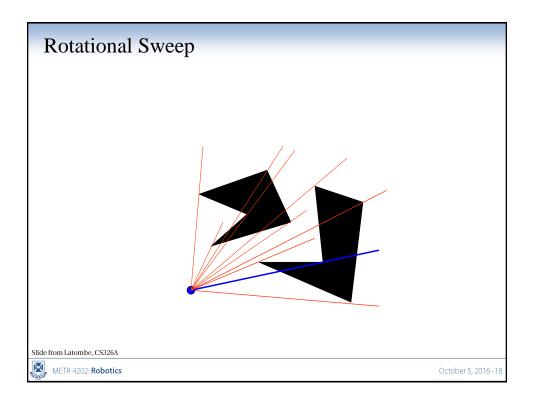


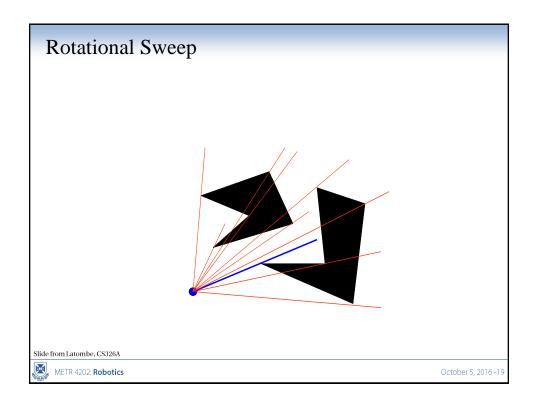


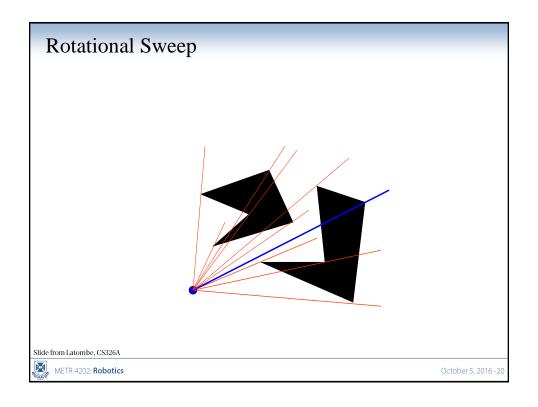










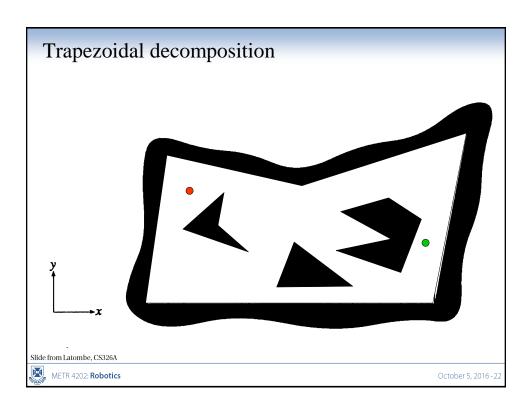


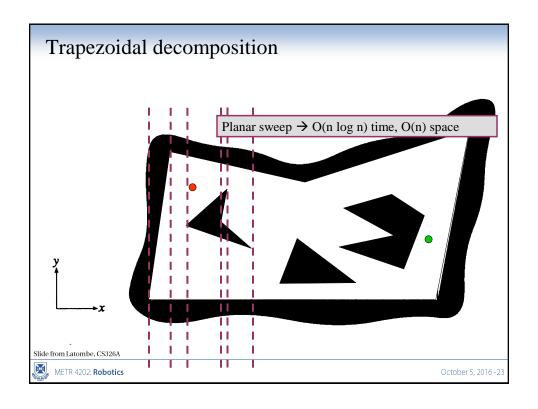
### II. Cell-Decomposition Methods

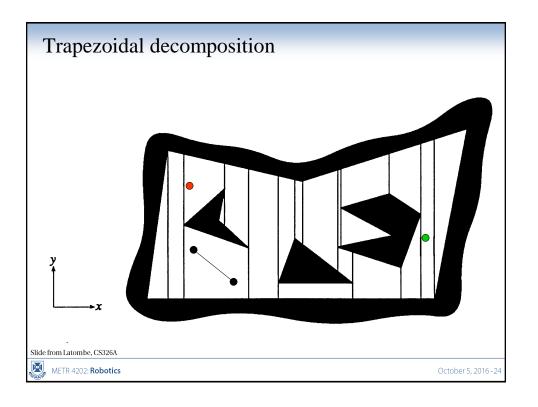
### Two classes of methods:

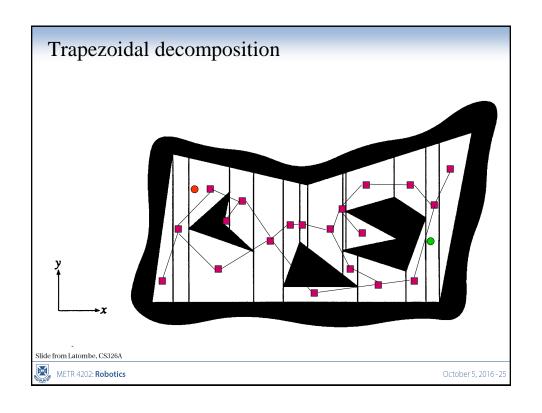
- Exact cell decomposition
  - The free space F is represented by a collection of nonoverlapping cells whose union is exactly F
  - Example: trapezoidal decomposition
- · Approximate cell decomposition
  - F is represented by a collection of non-overlapping cells whose union is contained in F Examples: quadtree, octree, 2n-tree

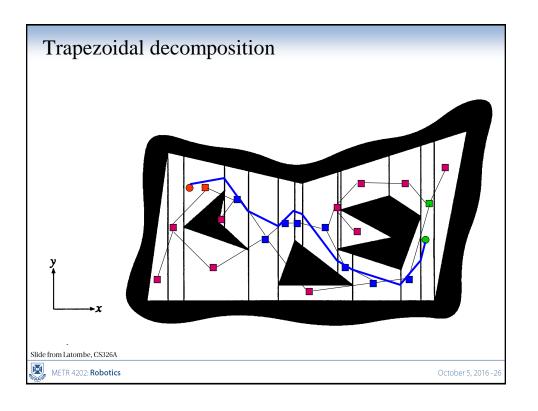


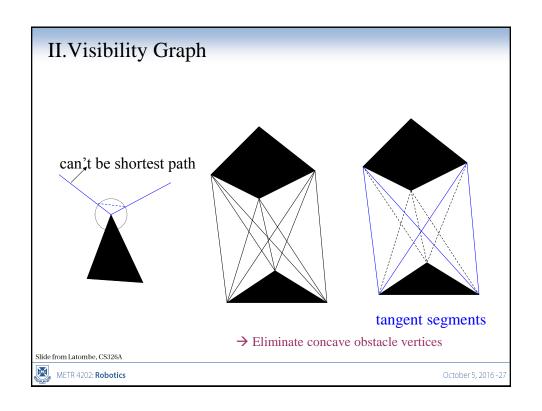


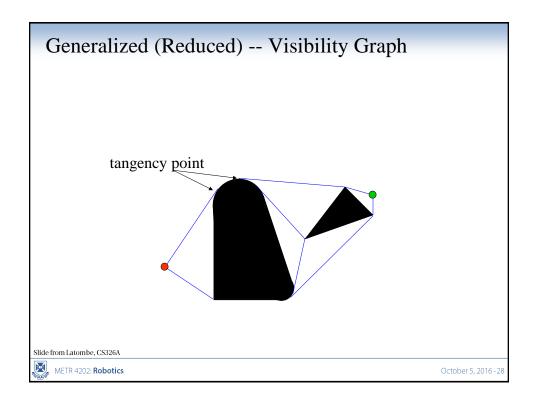


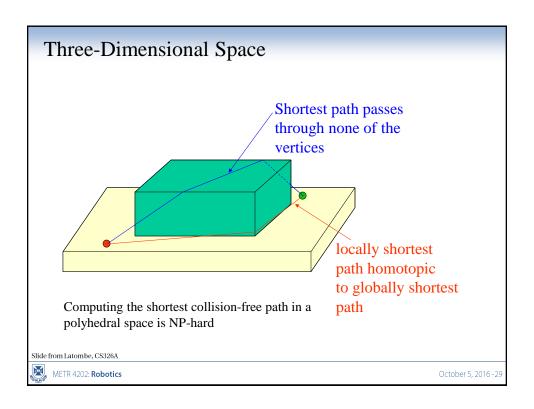






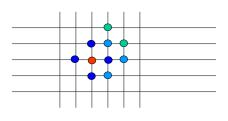






# Sketch of Grid Algorithm (with best-first search)

- Place regular grid G over space
- Search G using best-first search algorithm with potential as heuristic function



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### Simple Algorithm (for Visibility Graphs)

- Install all obstacles vertices in VG, plus the start and goal positions
- For every pair of nodes u, v in VG
   If segment(u,v) is an obstacle edge then

insert (u,v) into VG

else

for every obstacle edge e

if segment(u,v) intersects e then go up to segment

insert (u,v) into VG

• Search VG using A\*

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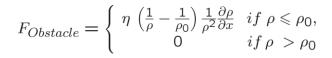
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### III. Potential Field Methods

• Approach initially proposed for real-time collision avoidance [Khatib, 86]

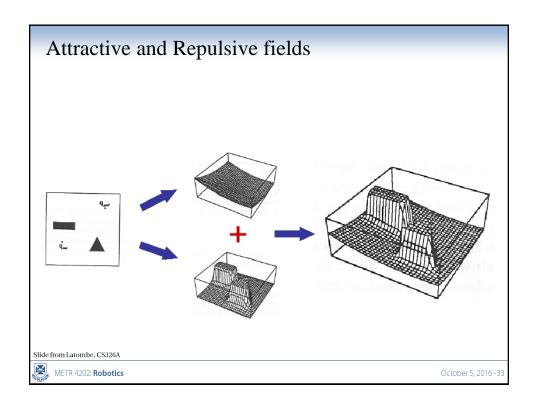
$$F_{Goal} = -k_p(x - x_{Goal})$$

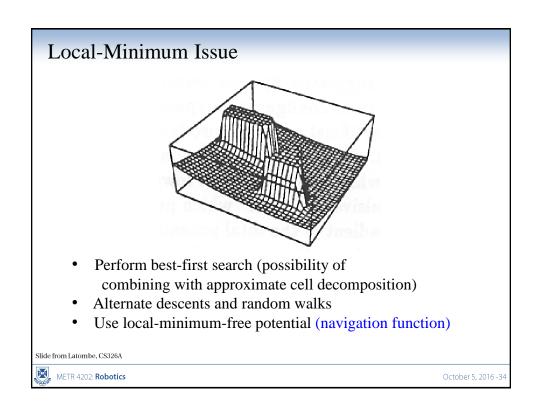


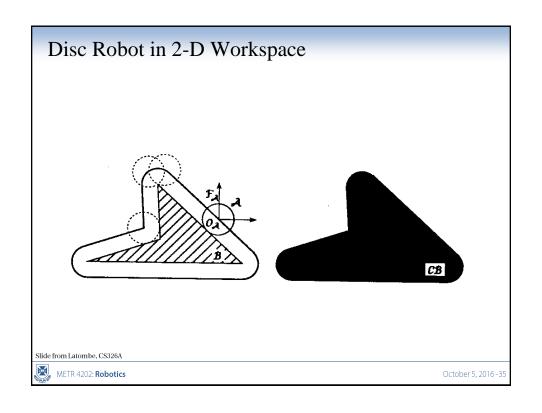
Coay Force Robot

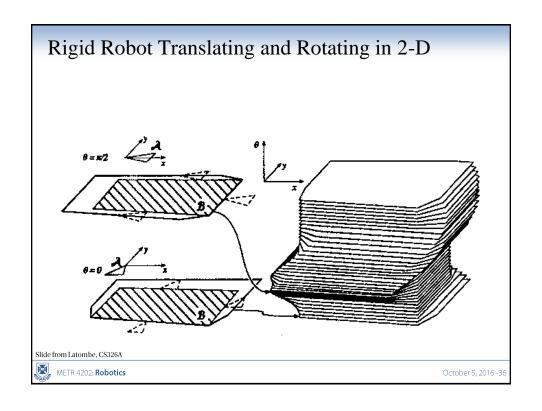
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### IV. Roadmap Methods

- · Visibility graph
- Voronoi diagram
- Silhouette
  First complete general method that applies to spaces of any dimension and is singly exponential in # of dimensions [Canny, 87]
- Probabilistic roadmaps (PRMS) and Rapidly-exploring Randomized Trees (RRTs)

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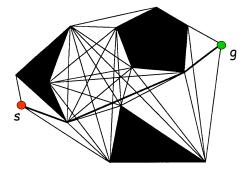


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### Roadmap Methods

· Visibility graph

Introduced in the Shakey project at SRI in the late 60s. Can produce shortest paths in 2-D configuration spaces



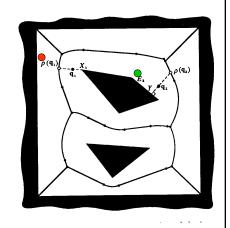
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# Roadmap Methods

Voronoi diagram
 Introduced by
 Computational
 Geometry researchers.
 Generate paths that maximizes clearance.

O(n log n) time O(n) space

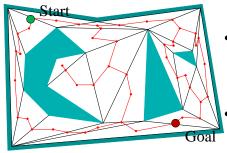


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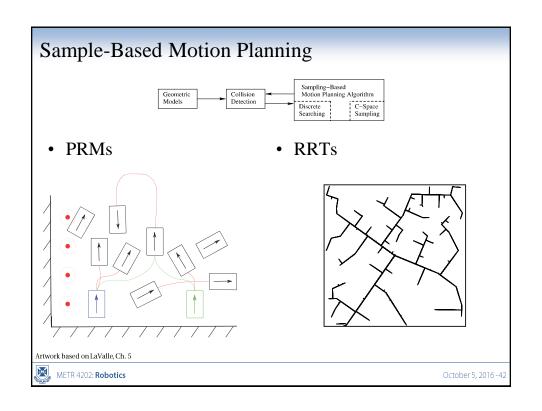
### Limits of Geometric Planning Methods

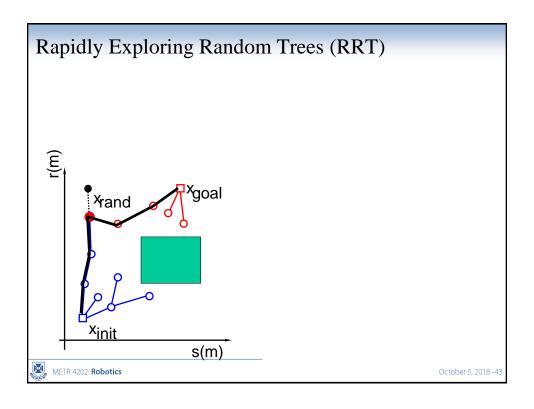


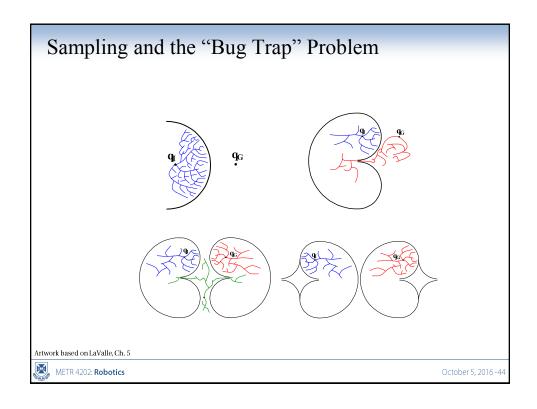
- How does this scale to high degrees of freedom?
- What about "dynamic constraints"?
- What about optimality?
- How to tie this to learning and optimization

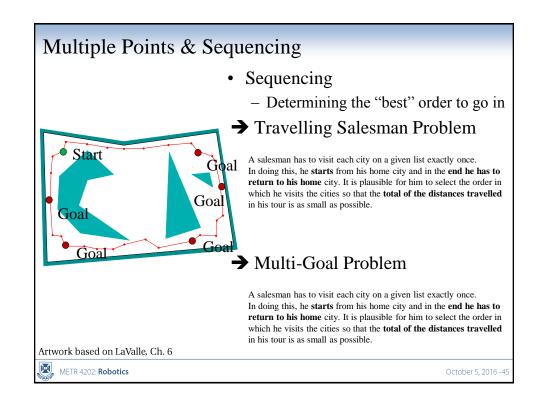
Artwork from LaValle, Ch. 6





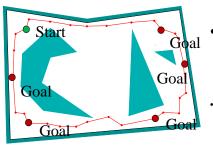






# Travelling Salesman Problem

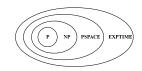
• Given a  $n \times n$  distance matrix  $\mathbf{C}=(c_{ii})$ 



• Minimize:

$$c(\pi) = \sum_{i=1}^{n} c_{i\pi(i)}$$

Note that this problem is NP-Hard



→ BUT, Special Cases are Well-Solvable!

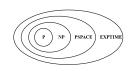
Artwork based on LaValle, Ch. 6



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### Travelling Salesman Problem [2]

• This problem is NP-Hard



→ BUT,
Special Cases are
Well-Solvable!

### For the Euclidean case

(where the points are on the 2D Euclidean plane):

- The shortest TSP tour does not intersect itself, and thus geometry makes the problem somewhat easier.
- If all cities lie on the boundary of a convex polygon, the optimal tour is a cyclic walk along the boundary of the polygon (in clockwise or counterclockwise direction).

### The k-line TSP

- The a special case where the cities lie on k parallel (or almost parallel) lines in the Euclidean plane.
- · EG: Fabrication of printed circuit boards
- Solvable in O(n<sup>3</sup>) time by Dynamic Programming (Rote's algorithm)

### The necklace TSP

 The special Euclidean TSP case where there exist n circles around the n cities such that every cycle intersects exactly two adjacent circles



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